

# Free $MV$ -algebras

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## Abstract

We describe  $n$ -generated free  $MV$ -algebras as  $MV$ -algebras having the lattice reduct which is a direct limit in the category of distributive lattices.

**Keywords:**  $MV$ -algebras, Lattices.

## 1 Preliminaries

We recall that an algebra  $A = (A; 0, 1, \oplus, \cdot, \sim), [1]$ , is said to be an  $MV$ -algebra iff it satisfies the following equations:

1.  $(x \oplus y) \oplus z = x \oplus (y \oplus z)$ ;
2.  $x \oplus y = y \oplus x$ ;
3.  $x \oplus 0 = x$ ;
4.  $x \oplus 1 = 1$ ;
5.  $\sim 0 = 1$ ;
6.  $\sim 1 = 0$ ;
7.  $x \cdot y = \sim(\sim x \oplus \sim y)$ ;
8.  $\sim(\sim x \oplus y) \oplus y = \sim(\sim y \oplus x) \oplus x$ .

Henceforth we shall write  $ab$  for  $a \cdot b$  and  $a^n$  for  $\underbrace{a \cdots a}_n$ , for given  $a, b \in A$ . Every  $MV$ -algebra has an underlying ordered structure defined by

$$x \leq y \text{ iff } \sim x \oplus y = 1.$$

$(A; \leq, 0, 1)$  is a bounded distributive lattice. Moreover, the following property holds in any  $MV$ -algebra:

$$xy \leq x \wedge y \leq x \vee y \leq x \oplus y.$$

The unit interval of real numbers  $[0, 1]$  endowed with the following operations:  $x \oplus y = \min(1, x + y)$ ,  $x \cdot y =$

$\max(0, x + y - 1)$ ,  $\sim x = 1 - x$ , becomes an  $MV$ -algebra. Let  $Q$  denote the set of rational numbers, for  $(0 \neq) m \in \omega$  we set  $S_m = (S_m; \oplus, \cdot, \sim, 0, 1)$ , where

$$S_m = \left\{ 0, \frac{1}{m}, \dots, \frac{m-1}{m}, 1 \right\}.$$

Recall that an algebra  $A \in \mathbf{K}$  is said to be a *free algebra* in a variety  $\mathbf{K}$ , if there exists a set  $A_0 \subset A$  such that  $A_0$  generates  $A$  and every mapping  $f$  from  $A_0$  to any algebra  $B \in \mathbf{K}$  is extended to a homomorphism  $h$  from  $A$  to  $B$ . In this case  $A_0$  is said to be *the set of free generators* of  $A$ . If the set of free generators is finite then  $A$  is said to be a *finitely generated free algebra*.

## 2 Representation of $n$ -generated $MV$ -algebras

Let us denote by  $DL$  the category, and the variety as well, of distributive lattices, named by  $DL$ -algebras in the sequel, and lattice homomorphisms.

Notice that any  $MV$ -algebra is bounded distributive lattice and the lattice operations  $\vee$  and  $\wedge$  are definable by means of  $MV$ -algebra operations  $\oplus$ ,  $\cdot$  and  $\sim$ . So, any  $MV$ -algebra is a  $DL$ -algebra with respect to the operations  $\vee, \wedge$ .

Let  $(A, \oplus, \cdot, \sim, 1, 0)$  be an  $MV$ -algebra. Define a functor  $\mathcal{D}$  transforming  $MV$ -algebra  $(A, \oplus, \cdot, \sim, 1, 0)$  into  $DL$ -algebra in the following way:

$$\mathcal{D}((A, \oplus, \cdot, \sim, 1, 0)) = (A, \vee, \wedge).$$

Then we have functor from the category of  $MV$ -algebras  $MV$  to the category  $DL$ .

Let  $F_{V_n}(m)$  be the  $m$ -generated free  $MV$ -algebra in the variety  $V_n = \mathcal{V}(\{S_1, \dots, S_n\})$ . We can consider any  $MV$ -algebra as (unbounded) distributive lattice (with lattice operations  $\vee$  and  $\wedge$ ). It is clear that  $F_{V_n}(m)$  is embedded into  $F_{V_{n+1}}(m)$  as  $MV$ -algebra and, consequently, as a lattice.

Let  $g_1^{(n)}, \dots, g_m^{(n)} \in F_{V_n}(m)$  be free generators of  $F_{V_n}(m)$ .

Let us define the injective mapping

$$\varepsilon_{n(n+1)} : \mathcal{D}(F_{V_n}(m)) \rightarrow \mathcal{D}(F_{V_{n+1}}(m))$$

for any  $n \in \omega/\{0\}$  in the following way:

$$g_1^{(n)} \mapsto g_1^{(n+1)}, \dots, g_m^{(n)} \mapsto g_m^{(n+1)} \text{ for any } n \in \omega/\{0\}.$$

It is clear that in the case when  $n = 1$ ,  $g_1^{(1)}, \dots, g_m^{(1)}$  generate  $F_{V_1}(m)$  ( $= D_1^{(m)} = \mathcal{D}(F_{V_1}(m))$ ) using only operations  $\vee, \wedge$  and  $\sim$ . Then  $\varepsilon_{12}(g_i^{(1)}) = g_i^{(2)}$ , for  $i = 1, \dots, m$ , and generate the algebra  $D_2^{(m)}$  inside  $F_{V_2}(m)$  by means of  $g_1^{(2)}, \dots, g_m^{(2)}$  using only operations  $\vee, \wedge$  and  $\sim$ . It easy to check that  $D_2^{(m)}$  is isomorphic to  $D_1^{(m)}$  ( $= \mathcal{D}(F_{V_1}(m))$ ).  $\varepsilon_{12}(x \vee y) = \varepsilon_{12}(x) \vee \varepsilon_{12}(y)$ ,  $\varepsilon_{12}(x \wedge y) = \varepsilon_{12}(x) \wedge \varepsilon_{12}(y)$ .

In particular, if  $m = 1$ , then  $\varepsilon_{12}(1) = g_1^{(2)} \vee \sim g_1^{(2)}$  ( $\neq 1$ ) and  $\varepsilon_{12}(0) = g_1^{(2)} \wedge \sim g_1^{(2)}$  ( $\neq 0$ ). Observe that  $\varepsilon_{12}(\mathcal{D}(F_{V_1}(1)))$  generates  $F_{V_1}(2)$  using  $MV$ -algebra polynomials containing operations  $\vee, \wedge, \sim$  and  $MV$ -algebra operation  $\oplus$  (or  $\cdot$ ) only one time. It is enough, since in  $V_2$  holds the axiom  $2x = 3x$  (and  $2x$  is a Boolean element). Indeed, by means of the generator  $g_1^{(2)}$  using  $\vee, \wedge, \sim$  and polynomials containing operation  $\oplus$  (or  $\cdot$ ) only one time, we obtain elements  $(0, 1, 1), (1, 0, 1), (1, 1, 0), (1, 0, 0), (0, 1, 0), (0, 0, 1)$  and, hence, any element of  $F_{V_1}(2)$  by means of  $\vee, \wedge$  and  $g_1^{(2)}, \sim g_1^{(2)}$ . Notice, that the embedding  $\varepsilon_{12}$  is not  $MV$ -algebra embedding, it is  $DL$ -algebra embedding. So, we define a functor  $\mathcal{L}_2$  transforming  $DL$ -algebra  $\mathcal{D}(F_{V_2}(1))$  into  $MV$ -algebra  $F_{V_2}(1)$ . Further,  $\varepsilon_{23}(g_1^{(2)}) = g_1^{(3)}$  and close  $\{g_1^{(3)}\}$  by means of operations  $\vee, \wedge, \sim$ . Then we obtain  $DL$ -algebra  $D_3^{(1)}$  which is isomorphic to  $\mathcal{D}(F_{V_1}(1))$ . Analogically, as in the previous case, close  $D_3^{(1)}$  using  $\vee, \wedge, \sim$  and polynomials containing operation  $\oplus$  (or  $\cdot$ ) only one time. Then we obtain  $DL$ -algebra  $\varepsilon_{23}(\mathcal{D}(F_{V_1}(2)))$  which is isomorphic to  $\mathcal{D}(F_{V_1}(2))$ . Further, close  $\varepsilon_{12}(\mathcal{D}(F_{V_1}(2)))$  using  $\vee, \wedge, \sim$  and polynomials containing operation  $\oplus$  (or  $\cdot$ ) only not more than on two places. Then we obtain the algebra  $\mathcal{D}(F_{V_1}(3))$ . So, we can define a functor  $\mathcal{L}_3$  transforming  $DL$ -algebra  $\mathcal{D}(F_{V_3}(1))$  into  $MV$ -algebra  $F_{V_3}(1)$ . And so on.

Analogically it is shown in general case when  $m > 1$ .

Further,  $\varepsilon_{12}(g_i^{(2)}) = g_i^{(3)}$ , for  $i = 1, \dots, m$ , and generate the algebra  $D_3^{(m)}$  inside  $F_{V_3}(m)$  by means of  $g_1^{(3)}, \dots, g_m^{(3)}$  using only operations  $\vee, \wedge$  and  $\sim$ . It easy to check that  $D_3^{(m)}$  is isomorphic to  $F_{V_1}(m)$  (as a lattice with involution). Moreover,  $D_3^{(m)}$  generates

$F_{V_3}(m)$ , if we use the  $MV$ -algebra operations  $\sim$  and  $\oplus$  only two times. In other words, for generating  $F_{V_3}(m)$  we use polynomials containing the operation  $\sim$  and  $\oplus$  only in not more than two places. It is enough, since in  $V_3$  holds the axiom  $3x = 4x$ .

Let us consider the direct system  $(\mathcal{D}(F_{V_n}(m)), \varepsilon_{n(n+1)})_{n \in \omega/\{0\}}$ . Let  $\mathcal{D}(F(m))$  be the direct limit of the direct system  $(\mathcal{D}(F_{V_n}(m)), \varepsilon_{n(n+1)})_{n \in \omega/\{0\}}$ . The elements of  $\mathcal{D}(F(m))$  are infinite sequences  $(a_i, a_{i+1}, a_{i+2}, \dots)$  for  $i \in \omega$  such that  $a_i \in \mathcal{D}(F_{V_i}(m))$  and  $a_i$  has no inverse image with respect to  $\varepsilon_{(i-1)i}$ , and, moreover,  $\varepsilon_{i_1(i_1+1)}(a_{i_1}) = a_{i_1+1}$  for any  $i_1 \geq i$ . Now we define the operation  $\oplus$  on  $\mathcal{D}(F(m))$  in the following way.

Let  $a = (a_i, a_{i+1}, a_{i+2}, \dots), b = (b_j, b_{j+1}, b_{j+2}, \dots)$  and  $a, b \in F(m)$ , and  $i \leq j$ . It is clear that  $a_i$  (respectively,  $b_j$ ) represented by a valuation of an  $MV$ -algebra polynomial containing operation  $\oplus$  only on  $i - 1$  (respectively,  $j - 1$ ) places. Then we consider elements  $a_{\max(i,j)}$  and  $b_{\max(i,j)}$ . It is clear that they belong to  $\mathcal{D}(F_{V_{\max(i,j)}}(m))$ . Choose the least number  $k \geq \max(i, j)$  such that  $a_{\max(i,j)} + b_{\max(i,j)} \in \mathcal{D}(F_{V_k}(m))$ . Then  $a \oplus b = (a_k + b_k, a_{k+1} + b_{k+1}, a_{k+2} + b_{k+2}, \dots)$ . Then we convert  $\mathcal{D}(F(m))$  into  $MV$ -algebra which we denote by  $F_{MV}(m)$ .

**Theorem 2.1.** *The  $MV$ -algebra  $F_{MV}(m)$  is  $m$ -generated free  $MV$ -algebra.*

## References

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